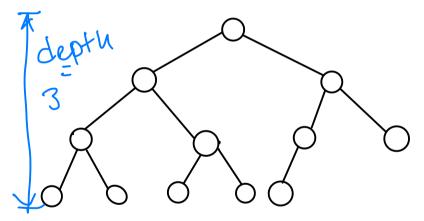
Lecture 14: Heaps Continued

- 1. Review Heap properties
- 2. Representing heaps w/ Arrays

Last Time

complete binary trees (CBTs)



Properties:

- 1. all nodes at depth = d-2 have 2 children
- 2. At most 7 node at depth d-1 has 1 child, and it is a left child
- 3. If v at depth d-1 has children and u is to the left (@ depth d), then u has 2 children
- 4. If v at depth d-1 has <2 children and w is to the right, then w has no children

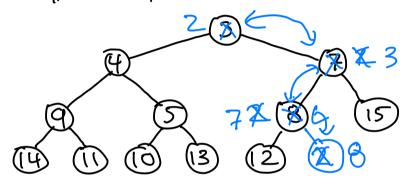
<u>Consequence</u>. If T is a complete binary tree then:

- 1. There is a unique location that a leaf can be added to result in a CBT
- 2. There is a unique leaf that can be removed to sesult in a CBT

Also last time: Heaps

A <u>heap</u> is a CBT in which each node stores a comparable element and satisfies:

Heap property the value stored at a node is no larger than the values stored by its childlen



Adding to a heap:

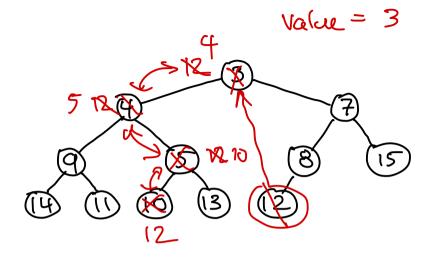
- 1. Add element at unique location to append a new leaf
- 2. "Bubble up": · v E new node
 - · while (v's val < v's parent's val)
 - swap values
 - set V = V. Parent

Example: add(2)

Removing min from a heap:

- 1. Store root val (to be returned)
- 2. Copy value from "last" leaf to root, and set as root value - right most leaf C clepth d
- 3. Remove leaf
- 4. "trickle down"
 - $\cdot v \leftarrow root$
 - · while (v's val > some child's val)
 - U = smaller child of V
 - swap u and v's vals
 - update ven
 - Example : remove Min()

<u>Results</u>: add and remove Min maintain heap property / CBT and can be performed with O(log n) compare/swap operations



Representing CBTs as Arrays

Since CBTs have predictable structure, we can represent them reatly as arrays:

- · root at index 0
- · left/right children @ indices 1,2
- · grand children C 3,4,5,6 (left to right)

We can access children/parents directly from the array:

- left child of index i is 2+i+1
- right child of index i is 2*i+2
- parent of index i is (i-1)/2

